

Algorithms: Complexity (Big Oh) of the codes

The RAM Model of Computation

Algorithms are an important and durable part of computer science because they can be studied in a machine/language independent way.

This is because we use the **RAM model of computation** for all our analysis.

- Each “simple” operation (+, -, =, if, call) takes 1 step.
- Loops and subroutine calls are *not* simple operations. They depend upon the size of the data and the contents of a subroutine. “Sort” is not a single step operation.
- Each memory access takes exactly 1 step.

We measure the run time of an algorithm by counting the number of steps.

This model is useful and accurate in the same sense as the flat-earth model (which *is* useful)!

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

```
for(int i=0; i<5; i++)  
    printf("%d", i);
```



$$c1 = 1 + 6 + 5 = 12$$



$$c2 = 5 * 1 = 5$$

$$\begin{aligned} \text{Total } c &= c1 + c2 = 12 + 5 = 17 \\ &= O(n^0) = O(1) \end{aligned}$$

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

```
for(int i=0; i<n; i++)
```

```
    printf("%d", i);
```



$$c1 = 1 + (n+1) + n = 2n + 2$$

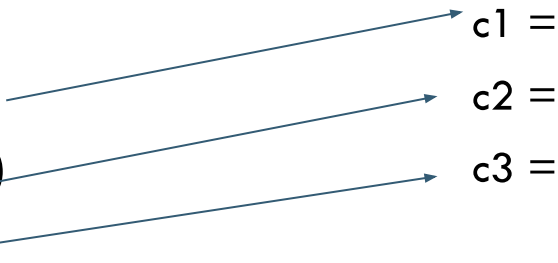
$$c2 = 1 * n = n$$

$$\text{Total } c = c1 + c2 = 3n + 2 = O(n)$$

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

```
for(int i=0; i<6; i++) c1 =  
    for(int k=0; k<5; k++) c2 =  
        printf("%d", i+k); c3 =
```

The diagram shows three lines of code. The first line is 'for(int i=0; i<6; i++)', with an arrow pointing from its end to the label 'c1 ='. The second line is ' for(int k=0; k<5; k++)', with an arrow pointing from its end to the label 'c2 ='. The third line is ' printf("%d", i+k);', with an arrow pointing from its end to the label 'c3 ='.

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

for(int i=0; i<6; i++)

for(int k=0; k<5; k++)

printf("%d", i+k);

$$c1 = 1 + (6+1) + 6 = 12$$

$$c2 = 6*1 + 6*(5+1) + 6*5 = 72$$

$$c3 = 6*5*1 = 30$$

$$\text{Total } c = c1 + c2 + c3$$

$$= 12 + 72 + 30 = 114$$

$$= O(n^0) = O(1)$$

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

```
for(int i=0; i<n; i++)
```

```
    for(int k=0; k<n; k++)
```

```
        printf("%d", i+k);
```

$$c1 = 1 + (n+1) + n = 2n + 2$$

$$c2 = n*1 + n*(n+1) + n*n = n*(2n + 2)$$

$$c3 = (n*n)*1 = n*n$$

$$\text{Total } c = c1 + c2 + c3 = (2n+2) + n*(2n+2) + n*n$$

$$= 3*n*n + 3*n + 2$$

$$= O(n*n)$$

Running Time/Cost of codes and Big Oh

- Assumption
 - Each operation will take same time

```
for(int i=0; i<m; i++)
```

```
    for(int k=0; k<n; k++)
```

```
        printf("%d", i+k);
```

$$c1 = 1 + (m+1) + m = 2m + 2$$

$$c2 = m*1 + m*(n+1) + m*n = m*(2n + 2)$$

$$c3 = (m*n)*1 = m*n$$

$$\text{Total } c = c1 + c2 + c3 = (2m+2) + m*(2n+2) + m*n$$

$$= 3*m*n + 3*m + 2$$

$$= O(m*n)$$

Asymptotic Analysis

High-level idea: Suppress **constant factors** and **lower-order terms**

too system-dependent

irrelevant for large inputs

Example: Equate $6n \log_2 n + 6$ with just $n \log n$.

Terminology: Running time is $O(n \log n)$

["big-Oh" of $n \log n$]

where $n =$ input size (e.g. length of input array).

Example: One Loop

Problem: Does array A contain the integer t ? Given A (array of length n) and t (an integer).

Algorithm 1

```
1: for  $i = 1$  to  $n$  do  
2:   if  $A[i] == t$  then  
3:     Return TRUE  
4: Return FALSE
```

Question: What is the running time?

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```

Question: What is the running time?

- A) $O(1)$ C) $O(n)$
B) $O(\log n)$ D) $O(n^2)$

Example: Two Loops

Given A, B (arrays of length n) and t (an integer). [Does A or B contain t ?]

Algorithm 2

```
1: for  $i = 1$  to  $n$  do
2:   if  $A[i] == t$  then
3:     Return TRUE
4: for  $i = 1$  to  $n$  do
5:   if  $B[i] == t$  then
6:     Return TRUE
7: Return FALSE
```

Question: What is the running time?

Example: Two Loops

Given A, B (arrays of length n) and t (an integer). [Does A or B contain t ?]

Algorithm 2

```
1: for  $i = 1$  to  $n$  do
2:   if  $A[i] == t$  then
3:     Return TRUE
4: for  $i = 1$  to  $n$  do
5:   if  $B[i] == t$  then
6:     Return TRUE
7: Return FALSE
```

→ Total $c = O(n)$

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Question: What is the running time?

- A) $O(1)$ C) $O(n)$ = $O(n) + O(n)$
B) $O(\log n)$ D) $O(n^2)$

Example: Two Nested Loops

Problem: Do arrays A, B have a number in common? Given arrays A, B of length n .

Algorithm 3

```
1: for  $i = 1$  to  $n$  do
2:   for  $j = 1$  to  $n$  do
3:     if  $A[i] == B[j]$  then
4:       Return TRUE
5: Return FALSE
```

Question: What is the running time?

Example: Two Nested Loops

Problem: Do arrays A, B have a number in common? Given arrays A, B of length n .

Algorithm 3

```
1: for  $i = 1$  to  $n$  do  
2:   for  $j = 1$  to  $n$  do  
3:     if  $A[i] == B[j]$  then  
4:       Return TRUE  
5: Return FALSE
```

Question: What is the running time?

- A) $O(1)$ C) $O(n)$
B) $O(\log n)$ D) $O(n^2)$

Example: Two Nested Loops (II)

Problem: Does array A have duplicate entries? Given arrays A of length n .

Algorithm 4

```
1: for  $i = 1$  to  $n$  do  
2:   for  $j = i+1$  to  $n$  do  
3:     if  $A[i] == A[j]$  then  
4:       Return TRUE  
5: Return FALSE
```

Question: What is the running time?

Example: Two Nested Loops (II)

Problem: Does array A have duplicate entries? Given arrays A of length n .

Algorithm 4

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1: for  $i = 1$  to  $n$  do
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```

Question: What is the running time?

- A) $O(1)$ C) $O(n)$
B) $O(\log n)$ D) $O(n^2)$

